

## Liveness:

```
spin -a file
gcc -o pan pan.c
pan -a -f or ./pan -a -f
spin -t -p -l -g -r -s file
```

## Spin arguments

-a	generate verifier and syntax check
-i	interactive simulation
-I	display Promela program after preprocessing
-nN	seed for random simulation
-t	guided simulation with trail
-tN	guided simulation with Nth trail
-uN	maximum number of steps is N
-f	translate an LTL formula into a never claim
-F	translate an LTL formula in a file into a never claim
-N	include never claim from a file
-l	display local variables
-g	display global variables
-p	display statements
-r	display receive events
-s	display send events

## Compile arguments

-DBFS	breadth-first search
-DNP	enable detection of non-progress cycles
-DSAFETY	optimize for safety
-DBITSTATE	bitstate hashing
-DCOLLAPSE	collapse compression
-DHC	hash-compact compression
-DMA=n	minimized DFA with maximum n bytes
-DMEMLIM=N	use up to N megabytes of memory

## Pan arguments

-a	find acceptance cycles
-f	weak fairness
-l	find non-progress cycles

-cN	stop after Nth error
-c0	report all errors
-e	create trails for all errors
-i	search for shortest path to error
-I	approximate search for shortest path to error
-mN	maximum search depth is N
-wN	$2^N$ hash table entries
-A	suppress reporting of assertion violations
-E	suppress reporting of invalid end states

## Caveats

- Expressions must be side-effect free.
- Local variable declarations always take effect at the beginning of a process.
- A true guard can always be selected; an else guard is selected only if all others are false.
- Macros and inline do *not* create a new scope.
- Place labels before an if or do, *not* before a guard.
- In an if or do statement, interleaving can occur between a guard and the following statement.
- Processes are activated and die in LIFO order.
- Atomic propositions in LTL formulas must be identifiers starting with lowercase letters and must be boolean variables or symbols for boolean-valued expressions.
- Arrays of bit or bool are stored in bytes.
- The type of a message field of a channel cannot be an array; it can be a typedef that contains an array.
- The functions empty and full cannot be negated.

## References

- G. J. Holzmann. *The Spin Model Checker: Primer and Reference Manual*, Addison-Wesley, 2004.  
<http://spinroot.com>.
- M. Ben-Ari. *Principles of the Spin Model Checker*, Springer, 2008.  
<http://www.springer.com/978-1-84628-769-5>.

# Spin Reference Card

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## Datatypes

bit	(1 bit)
bool	(1 bit)
byte	(8 bits unsigned)
short	(16* bits signed)
int	(32* bits signed)
unsigned	( $\leq$ 32* bits unsigned)

\* - for a 32-bit machine.

pid

chan

mtype = { name, name, ... } (8 bits)

typedef typename { sequence of declarations }

Declaration - type var [= initial value]

Default initial values are zero.

Array declaration - type var[N] [= initial value]

Array initial value assigned to all elements.

## Operators (descending precedence)

()	[]	.	
!	~	++	--
*	/	%	
+	-		
<<	>>		
<	=<	>	=>
==	!=		
&			
^			

```

|
&&
||
( ... -> ... : ... ) conditional expression
=

```

## Predefined

Constants - `true`, `false`

Variables (read-only except `_`):

- write-only hidden scratch variable
- `_nr_pr` - number of processes
- `_pid` - instantiation number of executing process
- `timeout` - no executable statements in the system?

## Preprocessor

```

#define name (arguments) string
#undef, #if, #ifdef, #ifndef, #else, #endif
#include "file name"
inline name (arguments) { ... }

```

## Statements

Assignment - `var = expression`, `var++`, `var--`  
`assert(expression)`

`printf`, `printm` - print to standard output  
`%c` (character), `%d` (decimal), `%e` (`mtype`),  
`%o` (octal), `%u` (unsigned), `%x` (hex)

`scanf` - read from standard input in simulation mode

`skip` - no operation

`break` - exit from innermost do loop

`goto` - jump to label

Label prefixes with a special meaning:

`accept` - accept cycle

`end` - valid end state

`progress` - non-progress cycle

`atomic { ... }` - execute without interleaving

`d_step { ... }` - execute deterministically (no jumping in or out; deterministic choice among true guards; only the first statement can block).

`{ ... } unless { ... }` - exception handling.

## Guarded commands

```

if :: guard -> statements :: ...
do :: guard -> statements :: ...
else guard - executed if all others are false.

```

## Processes

Declaration - `procname (parameters) { ... }`  
Activate with prefixes - `active` or `active[N]`  
Explicit process activation - `run procname (arguments)`  
Initial process - `init { ... }`  
Declaration suffixes:  
`priority` - set simulation priority  
`provided (e)` - executable only if expression e is true

## Channels

`chan ch = [ capacity ] of { type, type, ... }`

<code>ch ! args</code>	send
<code>ch !! args</code>	sorted send

<code>ch ? args</code>	receive and remove if <i>first</i> message matches
<code>ch ?? args</code>	receive and remove if <i>any</i> message matches
<code>ch ? &lt;args&gt;</code>	receive if <i>first</i> message matches
<code>ch ?? &lt;args&gt;</code>	receive if <i>any</i> message matches
<code>ch ? [args]</code>	poll <i>first</i> message (side-effect free)
<code>ch ?? [args]</code>	poll <i>any</i> message (side-effect free)

Matching in a receive statement: constants and `mtype` symbols must match; variables are assigned the values in the message; `eval(expression)` forces a match with the current value of the expression.

`len(ch)` - number of messages in a channel  
`empty(ch)` / `nempty(ch)` - is channel empty / not empty?  
`full(ch)` / `nfull(ch)` - is channel full / not full?

Channel use assertions:

<code>xr ch</code>	- channel ch is receive-only in this process
<code>xs ch</code>	- channel ch is send-only in this process

## Temporal logic

<code>!</code>	not
<code>&amp;&amp;</code>	and
<code>  </code>	or
<code>-&gt;</code>	implies
<code>&lt;-&gt;</code>	equivalent to
<code>[]</code>	always
<code>&lt;&gt;</code>	eventually
<code>X</code>	next
<code>U</code>	strong until
<code>V</code>	dual of U defined as <code>pVq &lt;-&gt; !(pU!q)</code>

## Remote references

Test the control state or the value of a variable:

<code>process-name @ label-name</code>
<code>procname-name [ expression ] @ label-name</code>
<code>process-name : label-name</code>
<code>procname-name [ expression ] : label-name</code>

## Never claim

`never { ... }`.

Predefined constructs that can only appear in a never claim:

<code>_last</code>	- last process to execute
<code>enabled(p)</code>	- is process enabled?
<code>np_</code>	- true if no process is at a progress label
<code>pc_value(p)</code>	- current control state of process
<code>remote references</code>	

See also `trace` and `notrace`.

## Variable declaration prefixes

<code>hidden</code>	- hide this variable from the system state
<code>local</code>	- a global variable is accessed only by one process
<code>show</code>	- track variable in Xspin message sequence charts

## Verification

Safety:

<code>spin -a file</code>
<code>gcc -DSAFETY -o pan pan.c</code>
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